

# Logical Constants and Focusing Disagreements

Greg Restall



THE UNIVERSITY OF  
MELBOURNE

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Motivation  
Defining Rules  
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# MOTIVATION

Logical vocabulary increases our expressive capacity.

Logical vocabulary is subject matter independent.

How do these fit together?

What distinctive things  
can we *do* with logical vocabulary?

# DEFINING RULES



$$X \vdash Y$$

Don't assert  $X$  and deny  $Y$ .

## Rules

$$\frac{X, A \vdash Y}{X \vdash \neg A, Y}$$

$$\frac{X, A, B \vdash Y}{X, A \wedge B \vdash Y}$$

$$\frac{X, A \vdash B, Y}{X \vdash A \rightarrow B, Y}$$

$$\frac{X \vdash A, B, Y}{X \vdash A \vee B, Y}$$

DISAGREEMENT

## Hinges

*... the questions that we raise and our doubts depend upon the fact that some propositions are exempt from doubt, are as it were like hinges on which those turn.*

*That is to say, it belongs to the logic of our scientific investigations that certain things are in deed not doubted. But it isn't that the situation is like this: We just can't investigate everything, and for that reason we are forced to rest content with assumption.*

*If I want the door to turn, the hinges must stay put.*

—Ludwig Wittgenstein, *On Certainty* §341–3

Logical vocabulary can function as fixed points,  
*around which* disagreement can take place.

## Example

$p \rightarrow q, p$ . Therefore,  $q$

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Accept  $p \rightarrow q$ , accept  $p$ , accept  $q$

Reject  $p \rightarrow q$ , accept  $p$ , accept  $q$

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~~Reject  $p \rightarrow q$ , reject  $p$ , reject  $q$~~

That looks like truth tables.

Can we look at this proof theoretically?

## Taking Positions

assert p	deny p
accept p	reject p
+p	-p

## Taking Positions

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assert p	<i>disclaim p</i>	deny p
accept p	<i>be open about p</i>	reject p
+p	?p	-p

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Think of  $?p$  as being *open* to both  $+p$  and  $-p$ .

(That works for the state, not the speech act.)

A POSITION is a collection of positively and negatively tagged statements.

- ▶ Two positions *disagree* when one contains  $+p$  and the other  $-p$  for some  $p$ .
- ▶ Two positions are *distinct* when one contains a positively or negatively tagged formula and the other doesn't.



Distinct positions on logically complex statements can be focused onto distinct positions on their constituents.

# BOOLEANS

## Example

Abelard:  $+(A \wedge B)$       Eloise:  $-(A \wedge B)$

## Example—Abelard

$$\frac{X, A, B \vdash Y}{\frac{\quad}{X, A \wedge B \vdash Y}}$$

$$A \wedge B \vdash A \quad A \wedge B \vdash B$$

Abelard cannot coherently add  $\neg A$  or  $\neg B$ .  
So he's not (*coherently*) open to  $\neg A$  and  $\neg B$ .  
He's committed (*implicitly*) to  $+A$  and to  $+B$ .

$$\frac{X, A, B \vdash Y}{X, A \wedge B \vdash Y}$$

$$A, B \vdash A \wedge B$$

Eloise cannot coherently add  $+A$  together with  $+B$ .  
So she's not (*coherently*) open to  $+A$  together with  $+B$ .

But she can be open to  $+A$ .  
And she can be open to  $+B$  too.

## Example (cont.)

<i>Abelard</i>	<i>Eloise</i>
	$-A, -B$
$+A, +B$	$+A, -B$
	$-A, +B$

## What about $\neg(A \wedge B)$

Abelard:  $\neg(A \wedge B)$       Eloise:  $\neg(A \wedge B)$

Then Abelard is *open to*  $+A$ ,  $+B$  and Eloise isn't, so there is still a difference concerning  $A$  and  $B$ .

## All the Booleans

$$\frac{X, A \vdash Y}{X \vdash \neg A, Y}$$

$$\frac{X, A, B \vdash Y}{X, A \wedge B \vdash Y}$$

$$\frac{X, A \vdash B, Y}{X \vdash A \rightarrow B, Y}$$

$$\frac{X \vdash A, B, Y}{X \vdash A \vee B, Y}$$

ONCE WE FIX THESE RULES, difference concerning negations, conjunctions, disjunctions, and conditionals can be focused into difference concerning their *constituents*.



(That's all well and good, but it still smells a lot like truth tables.)

# QUANTIFIERS

$$\frac{X \vdash A|_y^x, Y}{X \vdash \forall x A, Y}$$

$$\frac{X, A|_y^x \vdash Y}{X, \exists x A \vdash Y}$$

(where  $y$  is not free in  $X$  and  $Y$ )

Abelard:  $+\exists x A$       Eloise:  $-\exists x A$

$$\frac{X, A|_y^x \vdash Y}{X, \exists x A \vdash Y}$$

$$A|_y^x \vdash \exists x A$$

Eloise cannot coherently add  $+A|_y^x$  for any  $y$  at all.

So she's not (*coherently*) open to  $+A|_y^x$ .

So she's (*implicitly*) committed to  $-A|_y^x$  for any  $y$  whatever.

## Example—Abelard

$$\frac{X, A|_y^x \vdash Y}{X, \exists x A \vdash Y}$$

$$\frac{X, A|_y^x, \exists x A \vdash Y}{X, \exists x A \vdash Y}$$

Abelard *can* add  $+A|_y^x$  (for a new  $y$ ) at no cost to coherence.

*Why?* To simplify—to expose the structure of  $A$ .

If I assert  $\exists x A$ ,  
you can ask me to introduce a new term  $y$ ,  
and in that context I'm committed to  $+A|_y^x$ .

"You say  $\exists x A$ . Suppose that's right.  
Call it  $y$ . So  $A|_y^x \dots$ "

If I assert  $\exists x A$ ,  
you can ask me to introduce a new term  $y$ ,  
and in that context I'm committed to  $+A|_y^x$ .

“You say  $\exists x A$ . Suppose that's right.  
Call it  $y$ . So  $A|_y^x \dots$ ”



Differences over  $\exists x A$  (or  $\forall x A$ )  
can become differences over  $A$ ,  
at the cost of new vocabulary.

# MODALS

Differences over  $\diamond A$  (or  $\square A$ )  
can become differences over  $A$ ,  
at the cost of new *zones*.

If I assert  $\diamond A$ ,  
you can ask me to introduce a new zone  
and in that context I'm committed to  $+A$ .

"You say  $\diamond A$ . Suppose that's right.  
Consider if that obtained. Then,  $A \dots$ "

If I assert  $\diamond A$ ,  
you can ask me to introduce a new zone  
and in that context I'm committed to  $+A$ .

“You say  $\diamond A$ . Suppose that's right.  
Consider if that obtained. Then,  $A \dots$ ”

$$\frac{\mathcal{H}[A \vdash \quad | X \vdash Y]}{\mathcal{H}[X, \diamond A \vdash Y]}$$

## Example

$$\frac{\frac{p \vdash p \quad | \vdash \Diamond q}{p \vdash \quad | \vdash \Diamond p, \Diamond q} \quad \frac{q \vdash q \quad | \vdash \Diamond p}{q \vdash \quad | \vdash \Diamond p, \Diamond q}}{\frac{p \vee q \vdash \quad | \vdash \Diamond p, \Diamond q}{\Diamond(p \vee q) \vdash \Diamond p, \Diamond q}}}{\Diamond(p \vee q) \vdash \Diamond p \vee \Diamond q}$$

Of course...

... sense needs to be made of  $+A$  and  $-A$  in different zones.



# LIMITS

Enough of where this can work.

Where does it *fail*?

Differences concerning  $Fa$  needn't focus...

## Identity?

... including differences over  $a = b$  ...

... unless we are *very* generous  
about what counts as a constituent.

## Identity?

Abelard:  $+(a = b)$       Eloise:  $-(a = b)$

$a = b, Fa \vdash Fb$

Abelard is barred from  
 $+C|_a^x$  with  $-C|_b^x$ .

## But Eloise?

If Eloise denies  $a = b$ ,  
Abelard can ask her to introduce a new predicate  $X$   
and she is committed to  $+Xa$  and to  $-Xb$ .

“You deny  $a = b$ . Suppose they’re different. . .”



## But Eloise?

If Eloise denies  $a = b$ ,  
Abelard can ask her to introduce a new predicate  $X$   
and she is committed to  $+Xa$  and to  $-Xb$ .

“You deny  $a = b$ . Suppose they’re different...”

This is *stretching* the notion of a constituent (or a subformula)...

... but it's still wellfounded.

## Second Order Quantifiers?

Abelard:  $+\forall X A$       Eloise:  $-\forall X A$

$$\forall X A \vdash A|_{\lambda x.B}^X$$

Abelard is barred from  $\neg A|_{\lambda x.B}^X$ .

## But Eloise?

If Eloise denies  $\forall X A$ ,  
then Abelard can ask her  
to introduce a new predicate  $Y$   
and she is committed to  $\neg A|_Y^X$ .

Can sense be made of these predicate variables?

## Reductions

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$\neg A$	$\rightsquigarrow$	$A$
$A \wedge B$	$\rightsquigarrow$	$A, B$
$A \vee B$	$\rightsquigarrow$	$A, B$
$A \rightarrow B$	$\rightsquigarrow$	$A, B$

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$\forall x A$	$\rightsquigarrow$	$A _{x_0}^x, A _{x_1}^x, A _{x_2}^x, \dots$
$\exists x A$	$\rightsquigarrow$	$A _{x_0}^x, A _{x_1}^x, A _{x_2}^x, \dots$

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$a = b$	$\rightsquigarrow$	$X_0 a, X_0 b, X_1 a, X_1 b, X_2 a, X_2 b, \dots$
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$\forall X A$	$\rightsquigarrow$	$A _{X_0}^X, A _{X_1}^X, A _{X_2}^X, \dots$
$\exists X A$	$\rightsquigarrow$	$A _{X_0}^X, A _{X_1}^X, A _{X_2}^X, \dots$

---

(The formulas shrink. The transitive closure of  $\rightsquigarrow$  is well founded.)

We do not reduce  $\forall X A$  to *each* of its instances.  
That relation is not well founded.

$$\forall X X a > X a |_{\lambda x. \forall X X a}^X = \forall X X a$$

Differences over  $\forall X A$  resolve into  
differences over its free variable instances.



- ▶ GAIN: we can say new things with our new vocabulary.
- ▶ MODESTY: if we had *new* subject matter, we could differ on that, keeping our other commitments fixed.
  - What about *defined terms*? Suppose we treat '*bachelor*' as '*unmarried male*'? Can we disagree about whether something is a bachelor without disagreeing about '*unmarried*' and '*male*'?
  - What do we treat as a subformula in this case?

THANK YOU!